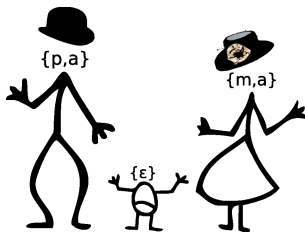


Set families with a forbidden subposet

Boris Bukh

August 2008



Introduction

$$[n] = \{1, \dots, n\} \quad 2^{[n]} = \{F : F \subset [n]\}$$

Definition

Set family is a collection of subsets of $[n]$. In symbols, $\mathcal{F} \subset 2^{[n]}$.

Theorem (Sperner'28)

Suppose $\mathcal{F} \subset 2^{[n]}$ is a set family such that for no distinct $F_1, F_2 \in \mathcal{F}$ the inclusion $F_1 \subset F_2$ holds. Then

$$|\mathcal{F}| \leq \binom{n}{\lfloor n/2 \rfloor}.$$

Equality is attained for $\mathcal{F} = \binom{[n]}{\lfloor n/2 \rfloor} = \{F \subset [n] : |F| = \lfloor n/2 \rfloor\}$.

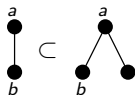
Reformulation

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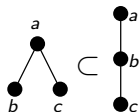
P_1 is a subposet of P_2 (written $P_1 \subset P_2$) if there is an injective $f: P_1 \rightarrow P_2$ such that

$$x <_{P_1} y \implies f(x) <_{P_2} f(y)$$

Examples:



and



Way of thinking

A set family is a poset under inclusion. For $F_1, F_2 \in \mathcal{F}$ set $F_1 \leq_{\mathcal{F}} F_2$ if $F_1 \subset F_2$.

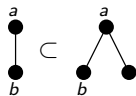
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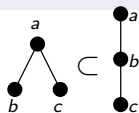
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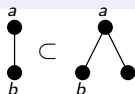
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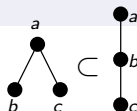
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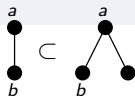
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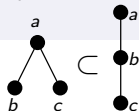
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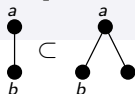
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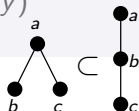
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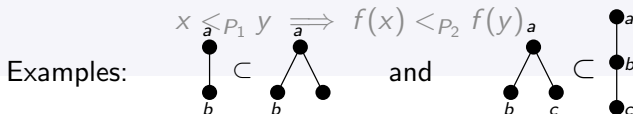
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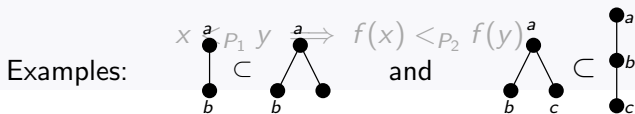
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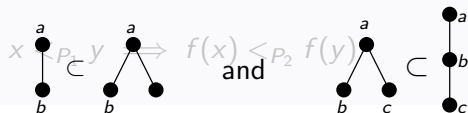
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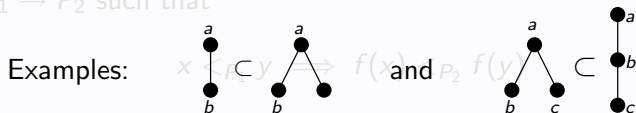
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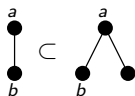


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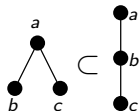
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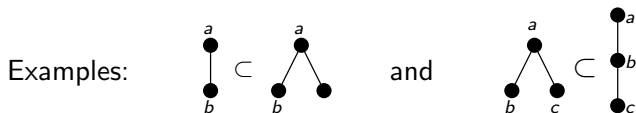
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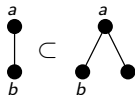


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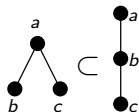
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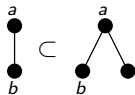


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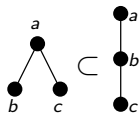
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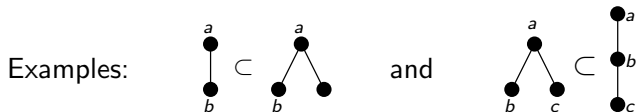
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Theorem (Sperner'28)

$$! \not\subset \mathcal{F} \implies |\mathcal{F}| \leq \binom{n}{\lfloor n/2 \rfloor}$$

Reformulation and the known results

Notation

Largest family in $[n]$ not containing poset P has size

$$\text{ex}(P, n) = \max_{\substack{P \notin \mathcal{F} \\ \mathcal{F} \subset 2^{[n]}}} |\mathcal{F}|$$

Sperner'28

$$\text{ex}(\uparrow, n) = \binom{n}{\lfloor n/2 \rfloor}$$

Erdős'45

$$\text{ex}(\uparrow, n) = \binom{n}{\lfloor n/2 \rfloor} + \binom{n}{\lfloor n/2 \rfloor + 1}$$

Katona–Tarján'83

$$\text{ex}(\vee, n) = \binom{n}{n/2} (1 + O(1/n))$$

Thanh'98

$$\text{ex}(\uparrow, n) = 2 \binom{n}{n/2} (1 + O(1/n))$$

De Bonis–Katona–Swanepoel'05

$$\text{ex}(\aleph, n) = \binom{n}{\lfloor n/2 \rfloor} + \binom{n}{\lfloor n/2 \rfloor + 1}$$

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Conjecture

For a fixed poset P

$$\text{ex}(P, n) = l(P) \binom{n}{n/2} (1 + O(1/n))$$

where $l(P)$ is the largest number of “middle” levels whose union contains no P .

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$$\text{ex}(\mathbb{N}, n) = \binom{n}{\lfloor n/2 \rfloor} + \binom{n}{\lfloor n/2 \rfloor + 1}$$

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where $l(P)$ is the largest number of “middle” levels whose union contains no P .

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Qualitatively explains all the previously known result including \aleph poset because $\aleph \subset \aleph'$.

For $h(P) = 2$ independently proved by Griggs and Lu. They also proved the conjecture for a large class of posets with $h(P) = 2$, whose Hasse diagram is not a tree.

First idea for $h(P) = 2$

Have: $|\mathcal{F}| \geq (1 + \varepsilon) \binom{n}{n/2}$, poset $P = \mathbb{N}$

Want: an embedding of P into poset (\mathcal{F}, \subset)

Idea

Treat poset (\mathcal{F}, \subset) as a graph, and embed tree into it.

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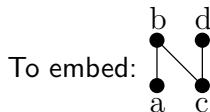
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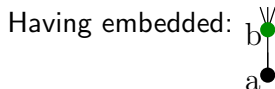
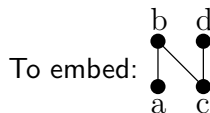
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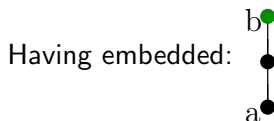
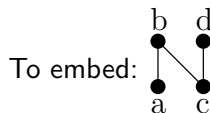
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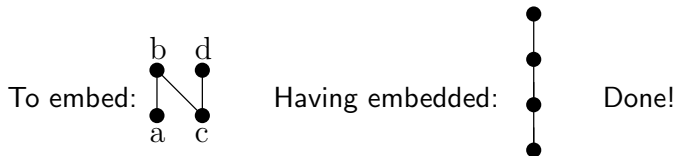
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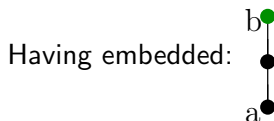
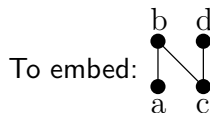
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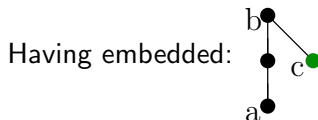
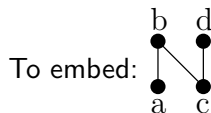
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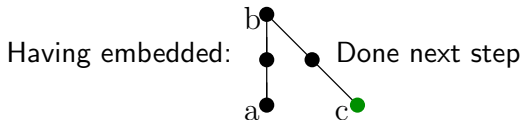
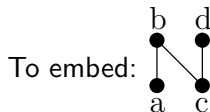
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For $h(P) \geq 3$ instead of graphs have to use $h(P)$ -uniform hypergraphs, and there is no good analogue of minimum degree.

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How to embed a tree into a graph of large average degree without using minimum degree?

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For every tree T there is a $d = d(T)$ such that every graph G of average degree $\geq d$ contains T .

Trouble with $h(P) \geq 3$

Difficulty

For $h(P) \geq 3$ instead of graphs have to use $h(P)$ -uniform hypergraphs, and there is no good analogue of minimum degree.

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Proof.

Induction on $|T|$. If $|T| = 1$, trivial. Let v be a leaf. Else let $T' = T \setminus \{v\}$, and $d(T) = 2d(T') + 4|T|$

- Let $V' = \{x \in V(G) : \deg(x) \geq d(T)/4\}$. Define $G' = G|_{V'}$
- Average degree of G' is at least $d(T)/2$.
- Find an embedding of T' into G' .
- Since $\deg(u) \geq |T|$ in G , extend the embedding.



□