

Small antipodal spherical codes

Boris Bukh

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Joint with Chris Cox

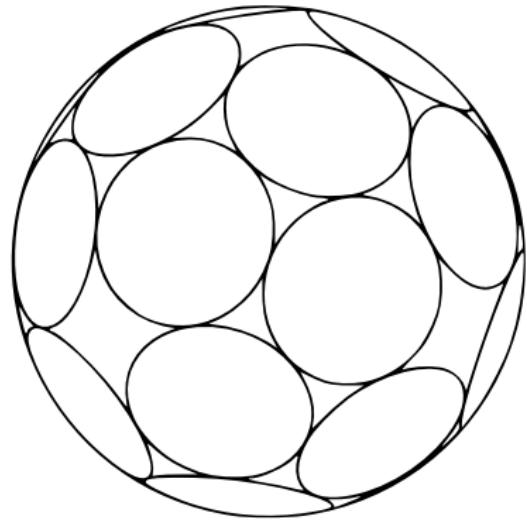
Spherical codes

Spherical code

unit vectors v_1, v_2, \dots, v_n in \mathbb{R}^d

Inner product

$$\max_{i \neq j} \langle v_i, v_j \rangle$$



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Wanted: large n and small inner product

Special case: binary codes, $v_i \in \{\pm 1/\sqrt{d}\}^d$

Antipodal codes

Antipodal code

vectors come in pairs $v, -v$

Inner product

$$\max_{i \neq j} |\langle v_i, v_j \rangle|$$

Basic problem

Unit vectors $v_1, v_2, \dots, v_{d+k} \in \mathbb{R}^d$,

$$f(d, k) = \min_{v_1, \dots, v_{d+k}} \max_{i \neq j} |\langle v_i, v_j \rangle|$$

Small antipodal codes

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Easy results:

$$f(d, 0) = 0 \quad \text{basis}$$

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$$f(d, 2) \lesssim \frac{2}{d} \quad \begin{aligned} &\text{simplex } \oplus \text{ simplex} \\ &\text{in } \mathbb{R}^{d/2} \oplus \mathbb{R}^{d/2} \end{aligned}$$

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$$f(d, k) \gtrsim \frac{\sqrt{k}}{d} \quad \begin{aligned} &\text{using } \text{tr}(A^2) \geq \text{tr}(A)^2 / \text{rk}(A) \\ &\text{(Welch bound; also LP bound)} \end{aligned}$$

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New results:

$$f(d, 2) \approx \frac{3/2}{d}$$

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$$f(d, k) \gtrsim \frac{c_k}{d} \quad \text{with } c_k = \frac{k(k+1)}{(k-1)\sqrt{k+2}+2}$$

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sharp for $k = 0, 1, 2, 3, 7, 23$

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sharp for $k = 0, 1, 2, 3, 7, 23$

$$f(d, k) \leq (1 + \varepsilon) \frac{\sqrt{k}}{d} \quad \text{where } \varepsilon \rightarrow 0 \text{ as } k \rightarrow \infty$$

Even sharper results for unit vectors in \mathbb{C}^d .

No linear programming

Isotropic measures

Isotropic measure μ on \mathbb{R}^k

$$\mathbb{E}_{x \sim \mu} |\langle x, v \rangle|^2 = \frac{1}{k} \|v\|^2 \quad \text{for all vectors } v$$

Key lemma

If μ is isotropic on \mathbb{R}^k , then

$$\mathbb{E}_{x,y \sim \mu} |\langle x, y \rangle| \leq \frac{(k-1)\sqrt{k+2} + 2}{k(k+1)}$$

Proof for the case $k = 1$

Unit vectors v_1, \dots, v_{d+1} such that $|\langle v_i, v_j \rangle| \leq \varepsilon$

Gram matrix $M = (\langle v_i, v_j \rangle)_{i,j}$

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$$M = \begin{pmatrix} 1 & [-\varepsilon, \varepsilon] & & \\ & 1 & & \\ & & 1 & \\ & & & \ddots \\ & [-\varepsilon, \varepsilon] & & 1 \\ & & & & 1 \end{pmatrix}$$

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Linear relation $0 = \sum \alpha_i r_i$ wlog $|\alpha_1| \geq |\alpha_2| \geq \dots \geq |\alpha_{d+1}|$

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Linear relation $0 = \sum \alpha_i r_i$ wlog $|\alpha_1| \geq |\alpha_2| \geq \dots \geq |\alpha_{d+1}|$

1st column $\alpha_1 = - \sum_{i \geq 2} \alpha_i r_{i1}$

$$\implies |\alpha_1| \leq \varepsilon \sum_{i \geq 2} |\alpha_i| \leq d\varepsilon |\alpha_1|$$

Proof for the case $k = 2$

$$\begin{pmatrix} 1 & & [-\varepsilon, \varepsilon] & & \\ & 1 & & & \\ & & 1 & & \\ & & & \ddots & \\ & & & & 1 \\ & & & & & 1 \end{pmatrix} \quad \begin{array}{ll} \text{Relation 1} & \text{Relation 2} \\ \alpha_1 & \beta_1 \\ \alpha_2 & \beta_2 \\ \alpha_3 & \beta_3 \\ \vdots & \vdots \end{array}$$

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		Relation 1	Relation 2	Relation
$\begin{pmatrix} 1 & & [-\varepsilon, \varepsilon] & \\ & 1 & & \\ & & 1 & \\ & & \ddots & \\ [-\varepsilon, \varepsilon] & & 1 & 1 \end{pmatrix}$		$\begin{matrix} \alpha_1 \\ \alpha_2 \\ \alpha_3 \\ \vdots \end{matrix}$	$\begin{matrix} \beta_1 \\ \beta_2 \\ \beta_3 \\ \vdots \end{matrix}$	$\begin{matrix} A\alpha_1 + B\beta_1 \\ A\alpha_2 + B\beta_2 \\ A\alpha_3 + B\beta_3 \\ \vdots \end{matrix}$

Seek (A, B) such that

$$|A\alpha_1 + B\beta_1| \geq \frac{3}{2} \mathbb{E}_i |A\alpha_i + B\beta_i|$$

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$$\begin{array}{ccc}
 & \text{Relation 1} & \text{Relation 2} \\
 \left(\begin{array}{cccc}
 1 & & [-\varepsilon, \varepsilon] & \\
 & 1 & & \\
 & & 1 & \\
 & & \ddots & \\
 & & & 1 \\
 [-\varepsilon, \varepsilon] & & & 1
 \end{array} \right) &
 \begin{array}{ccc}
 \alpha_1 & \beta_1 & A\alpha_1 + B\beta_1 \\
 \alpha_2 & \beta_2 & A\alpha_2 + B\beta_2 \\
 \alpha_3 & \beta_3 & A\alpha_3 + B\beta_3 \\
 \vdots & \vdots & \vdots
 \end{array} &
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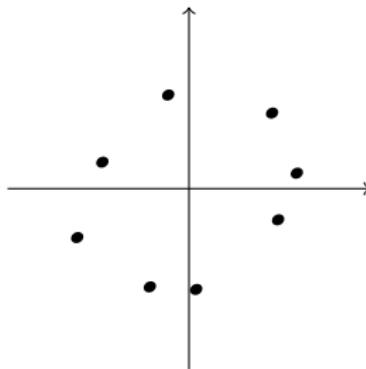
Points $p_i = (\alpha_i, \beta_i)$. Seek $L(x, y) = Ax + By$ such that

$$|Lp_1| \geq \frac{3}{2} \mathbb{E}_i |Lp_i|$$

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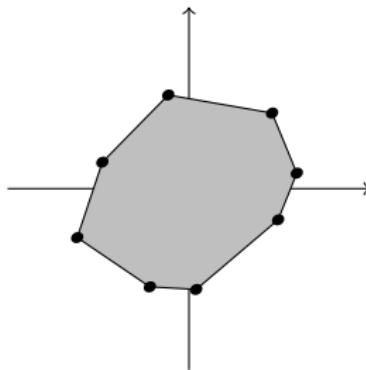
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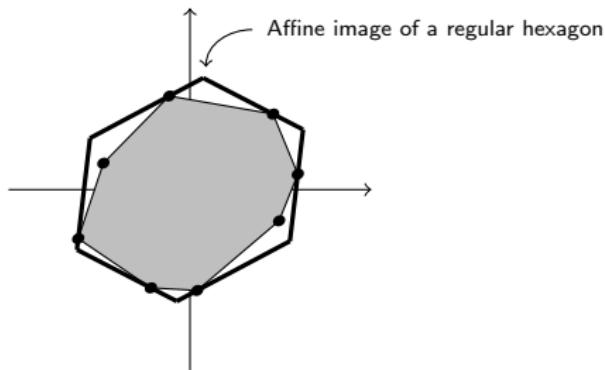
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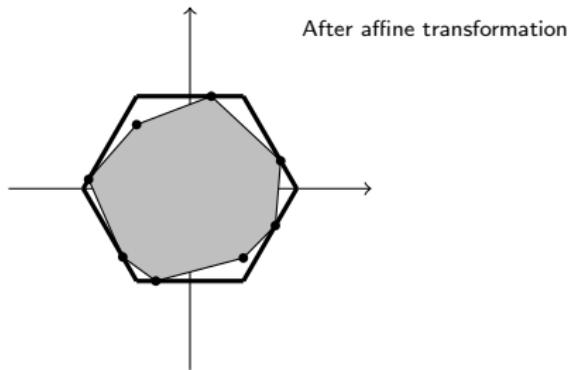
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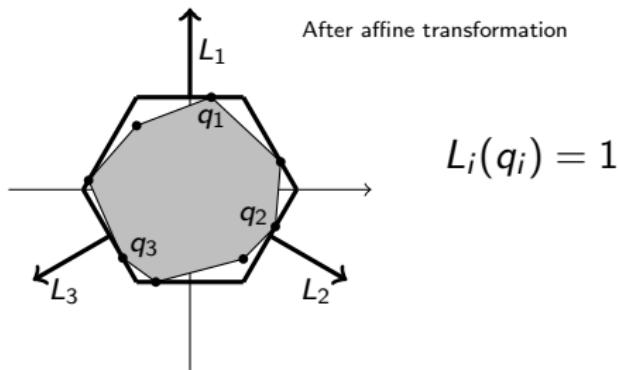
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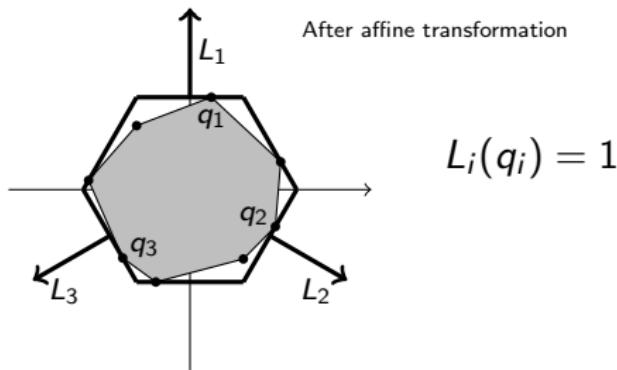
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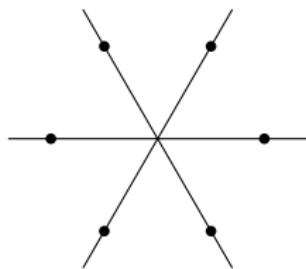
$$|Lp_1| \geq \frac{3}{2} \mathbb{E}_i |Lp_i|$$



$$\begin{aligned}\forall p \in \bullet \quad & |L_1(p)| + |L_2(p)| + |L_3(p)| \leq 2 \\ \implies \mathbb{E}_i |L_1(p_i)| + |L_2(p_i)| + |L_3(p_i)| & \leq 2 \\ \implies \mathbb{E}_i |L_j(p_i)| & \leq 2/3 \text{ for some } j\end{aligned}$$

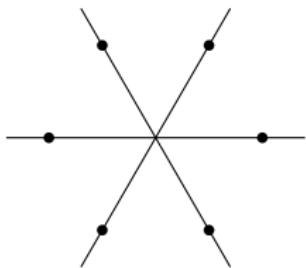
Construction for special k

$k = 2$:



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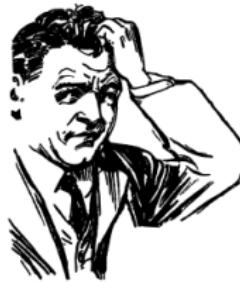
Equiangular lines

every angle is the same

In \mathbb{R}^k at most $\binom{k+1}{2}$ equiangular lines
equality for $k = 0, 1, 2, 3, 7, 23$

Problems

- ℓ^p norm of $|\langle v_i, v_j \rangle|_{i \neq j}$ (this result is $p = \infty$)
 $\implies \ell^q$ norm of isotropic measures for $\frac{1}{p} + \frac{1}{q} = 1$
- Non-antipodal codes ($2d+k$ points)

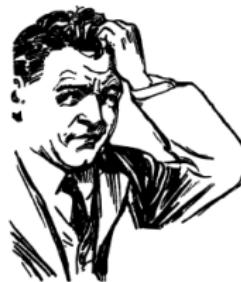


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10hrs ago: Alexey Glazyrin announced a solution

- Non-antipodal codes ($2d+k$ points)
 $\Omega(1/d)$ bound (w/Igor Balla)



Well, you asked for it!

Lemma

If μ is isotropic on \mathbb{R}^k , then

$$\mathbb{E}_{x,y \sim \mu} |\langle x, y \rangle| \leq \frac{(k-1)\sqrt{k+2} + 2}{k(k+1)}$$

Key quantity:

$$\mathbb{E}_{x,y} \left(\frac{|\langle x, y \rangle|}{\sqrt{\|x\| \|y\|}} - \beta \sqrt{\|x\| \|y\|} \right)^2$$

Upper bound:

- 1) Expand
- 2) Cauchy-Schwarz

Lower bound:

- 1) Cauchy-Schwarz
- 2) $\text{tr}(A^2) \geq \text{tr}(A)^2 / \text{rk}(A)$