

# Geometric combinatorics: supplementary notes 1\*

**Theorem 1** (Jung's theorem). *If  $X \subset \mathbb{R}^d$  is a finite set such that  $\|x - y\| \leq 1$  for all  $x, y \in X$ , then there is a point  $p \in \mathbb{R}^d$  such that*

$$\|p - x\| \leq \frac{1}{\sqrt{2}} \left( \frac{d}{d+1} \right)^{1/2}$$

for all  $x \in P$ .

*Proof.* Since the goal is to prove that the closed balls of radius  $\frac{1}{\sqrt{2}} \left( \frac{d}{d+1} \right)^{1/2}$  centred at the points of  $X$  have a non-empty intersection, by Helly's theorem it suffices to treat the case  $|X| \leq d + 1$ .

Without loss of generality, we can assume that the centre of the smallest ball containing  $X$  is the origin, and the radius of this ball is  $r$ . Let  $x_1, \dots, x_m$  be all the points of  $X$  that are at distance exactly  $r$  from 0.

The minimality of  $r$  implies that  $0 \in \text{conv}\{x_1, \dots, x_m\}$ . Indeed, if  $0 \notin \text{conv}\{x_1, \dots, x_m\}$ , then there is a vector  $v \in \mathbb{R}^d$  such that  $\langle v, x_i \rangle > 0$  for all  $i = 1, \dots, m$ , and the translated set  $P - \varepsilon v$  is closer to the origin than  $P$  was, provided  $\varepsilon$  is small enough. Since  $0 \in \text{conv}\{x_1, \dots, x_m\}$ , there are  $\alpha_1, \dots, \alpha_k$  such that

$$\begin{aligned} \alpha_1 x_1 + \dots + \alpha_m x_m &= 0, \\ \alpha_1 + \dots + \alpha_k &= 1. \end{aligned}$$

Hence

$$\begin{aligned} \sum_{i \neq j} \alpha_i &\geq \sum \alpha_i \|x_i - x_j\|^2 = \sum_i \alpha_i (2r^2 - 2\langle x_i, x_j \rangle) \\ &= 2r^2 - \left\langle \sum_i \alpha_i x_i, x_j \right\rangle = 2r^2. \end{aligned}$$

Averaging over all  $j = 1, \dots, m$  we obtain  $\frac{m-1}{m} \sum_i \alpha_i \geq 2r^2$ . The result follows from  $m \leq d + 1$ .  $\square$

It is easy to infer from the proof above that the theorem is sharp only if  $X$  contains all the  $d + 1$  vertices of a regular simplex with side-length 1.

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\*These notes are from <http://www.borisbukh.org/GeoCombEaster10/supnotes1.pdf>.

**Lemma 2.** *If  $C \subset \mathbb{R}^d$  is a non-empty compact set, there is a lexicographically minimal point of  $C$ .*

*Proof.* The proof is by induction on  $d$ . The base case  $d = 0$  is trivial. If  $d \geq 1$ , let  $x_{\min} = \inf\{x_1 : (x_1, \dots, x_d) \in C\}$ . Since  $(x_1, \dots, x_d) \mapsto x_1$  is continuous, the value  $x_{\min}$  is attained, and  $C' = \{y \in \mathbb{R}^{d-1} : \{x_{\min}\} \times y\}$  is non-empty. Since  $C'$  is compact subset of  $\mathbb{R}^{d-1}$ , the lemma follows by induction.  $\square$